

GENERAL DECOMPOSITION THEOREMS FOR m -CONVEX SETS IN THE PLANE

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ABSTRACT

A set S in R^d is said to be m -convex, $m \geq 2$, if and only if for every m distinct points in S , at least one of the line segments determined by these points lies in S . Clearly any union of $m - 1$ convex sets is m -convex, yet the converse is false and has inspired some interesting mathematical questions: Under what conditions will an m -convex set be decomposable into $m - 1$ convex sets? And for every $m \geq 2$, does there exist a $\sigma(m)$ such that every m -convex set is a union of $\sigma(m)$ convex sets? Pathological examples convince the reader to restrict his attention to closed sets of dimension ≤ 3 , and this paper provides answers to the questions above for closed subsets of the plane.

If S is a closed m -convex set in the plane, $m \geq 2$, the first question may be answered in one way by the following result: If there is some line H supporting S at a point p in the kernel of S , then S is a union of $m - 1$ convex sets. Using this result, it is possible to prove several decomposition theorems for S under varying conditions. Finally, an answer to the second question is given: If $m \geq 3$, then S is a union of $(m - 1)^3 2^{m-3}$ or fewer convex sets.

1. Introduction

Let S be a subset of R^d . The set S is said to be m -convex, $m \geq 2$, if and only if for every m distinct points in S , at least one of the line segments determined by these points lies in S . A point x in S is said to be a *point of local convexity* of S if and only if there is some neighborhood N of x such that if $y, z \in S \cap N$, then $[y, z] \subseteq S$. If S fails to be locally convex at some point q in S , then q is called a *point of local nonconvexity* (Inc point) of S . The following familiar terminology will be used: For x, y in S , we say x sees y via S if and only if the corresponding segment $[x, y]$ lies in S . Points x_1, \dots, x_n in S are *visually independent via S* if and only if for $1 \leq i < j \leq n$, x_i does not see x_j via S . Throughout the paper, $\text{conv } S$, $\text{ker } S$, $\text{bdry } S$, and $\text{cl } S$ will be used to denote the convex hull of S , the kernel of S , the boundary of S , and the closure of S , respectively. For convenience, Q will represent the set of Inc points of S .

Received March 15, 1975 and in revised form April 2, 1976

Several interesting decomposition theorems have been obtained for closed 3-convex sets in the plane. Valentine [9] has proved that a closed planar 3-convex set S may be written as a union of three or fewer convex sets. If, in addition, S is bounded and has some point of local convexity in $\text{bdry } S \cap \ker S$, then by a result of Stamey and Marr [7], S is a union of two convex sets.

For m -convex sets, we have the following analogue (Breen, [1]): For S a closed planar m -convex set with $\text{inc } Q$ points in Q , if $\text{conv } Q \subseteq S$ and $[(\text{bdry } S) \cap (\ker S)] \sim Q \neq \emptyset$, then S is a union of $m - 1$ closed convex sets. However, few other results have been obtained for the general case. Examples by Kay and Guay [5, Example 4] show that such a generalization must require an unpleasantly large number of convex sets, and it was only recently proved by Eggleston [2] that a compact planar m -convex set is expressible as a finite union of convex sets. Here we establish actual bounds for Eggleston's theorem using entirely different methods of proof. Several smaller bounds are obtained in case $\ker S \neq \emptyset$ or $\text{conv } Q \subseteq S$. Also, for $m = 4$, the bound of 6 is established.

2. The case for $\ker S \neq \emptyset$

Theorem 1 employs a basic construction introduced in [1] to generalize a result of that paper. The following theorem by Lawrence, Hare and Kenelly [6, Theorem 2] will be useful throughout the proof.

LAWRENCE, HARE, KENELLY THEOREM. *Let T be a subset of a linear space such that each finite subset $F \subseteq T$ has a k -partition $\{F_1, \dots, F_k\}$, where $\text{conv } F_i \subseteq T$, $1 \leq i \leq k$. Then T is a union of k convex sets.*

THEOREM 1. *Let S be a closed m -convex set in the plane, $m \geq 2$. Let $p \in \ker S \neq \emptyset$, and for H some line containing p , assume that $S \subseteq \text{cl } H_1$ (where H_1 is one of the open halfspaces determined by H). Then S is the union of $m - 1$ convex sets. The result is best possible for all m .*

PROOF. The proof will require several steps: First we show that we may assume S to be bounded and Q to be finite (Lemma 1), with $p \notin Q$ (Lemma 2). Then we consider the collection of rays at p consisting of rays of the form $R(p, q_i)$ for q_i in Q , together with rays R_1 and R_2 , where $R_1 \cup R_2 = H$. Order the rays appropriately. Each pair of consecutive rays will define a convex subset of S called a wedge, and we decompose S by defining $m - 1$ collections of wedges (Lemmas 3 and 4).

We begin by noticing that we may restrict our attention to the case in which S is bounded: For F any finite subset of S , F lies in some compact disk B , $B \cap S$

satisfies the hypothesis above, and by the Lawrence, Hare, Kenelly Theorem it suffices to prove the result for $B \cap S$. Therefore we shall assume that S is bounded.

LEMMA 1. *To any finite subset $F = \{x_i : 1 \leq i \leq n\}$ of S there corresponds an m -convex set T having finitely many lnc points, with $F \subseteq T \subseteq S$. Hence we may assume that S has finitely many lnc points.*

PROOF OF LEMMA 1. Let R_1, R_2 be closed rays at p , with $R_1 \cup R_2 = H$. Consider the family \mathcal{R} of rays consisting of R_1, R_2 together with rays of the form $R(p, q)$ emanating from p through q for some q in Q . It is not hard to show that for R in \mathcal{R} , R contains at most two members of Q . Any two (not necessarily distinct) rays in \mathcal{R} bound a closed subset of S , and we let \mathcal{W} denote the collection of all these closed regions. Moreover, since Q is closed, to every point x of S there corresponds a minimal member A of \mathcal{W} which contains x .

Now let $F = \{x_i : 1 \leq i \leq n\}$ be a finite subset of S . To each x_i there corresponds a minimal member A_i of \mathcal{W} which contains x_i . Each lnc point of S in A_i must lie in one of the boundary rays of A_i ; hence A_i contains at most four members of $Q \sim \{p\}$.

By arguments given in [1], we may assume that no A_i is a segment and also that $A_i = \text{cl}(\text{int } A_i)$. (In case $p \in Q$, the A_i sets are not necessarily convex; however, standard arguments show each component of $A_i \sim \{p\}$ to be convex.) Now order the rays associated with the A_i sets in a clockwise direction from R_1 (for an appropriate labeling of R_1, R_2). This in turn induces an order among the A_i sets, and we may relabel the A_i and corresponding x_i so that for $i < j$, the rays defining A_i precede the rays defining A_j in our clockwise ordering. Since $A_i = \text{cl}(\text{int } A_i)$, then each wedge is associated with at most two lnc points from $Q \sim \{p\}$, denoted q_i, q'_i . (In case A_i is associated with one lnc point in $Q \sim \{p\}$, then A_i must be bounded by R_1 or R_2 , and then we let q_1 be the last point of S on R_1 , q'_n the last point of S on R_2 .)

By the Lawrence, Hare, Kenelly Theorem, we may assume that each A_i has polygonal boundary, and we may select p_i, p'_i so that $[p_i, q_i]$ and $[p'_i, q'_i]$ lie in bdry A_i .

For $1 \leq i \leq n - 1$, let B_i denote the union of all segments $[x, y]$, where $[x, y] \subseteq S$, $x \in [p, q'_i]$, $y \in [p, q_{i+1}]$. We assert that $(\text{conv } B_i) \sim B_i$ is convex and $(\text{bdry conv } B_i) \sim B_i$ is polygonal: For s, t in $(\text{conv } B_i) \sim B_i$, if $[s, t]$ were not in $(\text{conv } B_i) \sim B_i$, then (s, t) would contain some point u in B_i , and for some x in $[p, q'_i]$, y in $[p, q_{i+1}]$, $u \in (x, y) \subseteq S$. But then one of s, t would lie on the p side of

the line $L(x, y)$, clearly impossible since $\text{conv}\{p, x, y\} \subseteq B_i$. Therefore $[s, t] \subseteq (\text{conv } B_i) \sim B_i$, and the set is convex. Since S is closed, $(\text{bdry conv } B_i) \sim [q'_i, q_{i+1}] \subseteq S$, and since S is m -convex, it is easy to see that $(\text{bdry conv } B_i) \sim [q'_i, q_{i+1}]$ is polygonal and consists of at most $m - 1$ segments. Hence the set B_i has at most $m - 2$ lnc points.

Define $T = \cup\{A_i \cup B_i \cup A_n : 1 \leq i \leq n - 1\}$. We assert that T is closed and m -convex, that the set Q_T of lnc points of T is finite, and that $p \in \ker T$. For any m -point subset of T , at least one of the corresponding segments, say $[v, w]$, is in S . We will show that $[v, w] \subseteq T$. In case $v, w \in A_i$ for some i , the result is trivial. If $v, w \in B_i$, then if (v, w) contained some point not in B_i , (v, w) would contain two boundary points of $(\text{conv } B_i) \sim B_i$, neither in (q'_i, q_{i+1}) . Thus the line $L(v, w)$ would intersect both $[p, q'_i]$ and $[p, q_{i+1}]$, forcing $[v, w]$ to lie in B_i , a contradiction. Hence $[v, w] \subseteq B_i$.

In case v, w do not lie in the same A_i or B_i set, then since $\text{conv}\{p, v, w\} \subseteq S$, no point of Q can lie interior to $\text{conv}\{p, v, w\}$. Therefore $[v, w]$ must intersect each $[p, q_i]$ and $[p, q'_i]$ between $R(p, v)$ and $R(p, w)$, and $[v, w]$ can be written as a finite union of segments in S , each having end points in some A_i or B_i set. Therefore, by previous remarks, each of these segments is in T , and $[v, w] \subseteq T$. It is clear that Q_T is finite, since at most two lnc points are contributed by each $A_i \sim \{p\}$, and at most $m - 2$ by each B_i .

Returning to a consideration of what will be needed to prove the theorem, in view of the Lawrence, Hare, Kenelly Theorem it suffices to prove that the set T just constructed is a union of $m - 1$ convex sets. Since clearly $p \in \ker T$, $T \subseteq \text{cl } H_1$, and Q_T is finite, it is therefore sufficient to prove the theorem under the assumption that S is bounded and Q is finite.

LEMMA 2. *We may assume that p is not an lnc point for S .*

PROOF OF LEMMA 2. By Lemma 1, we may assume that the set Q of lnc points of S is finite, so we may select some convex neighborhood N of p such that $N \cap Q \sim \{p\}$ is empty. Using standard arguments, it is easy to show that each component of $N \cap S \sim \{p\}$ has convex closure, and by remarks in [1], we may assume that no such component is a segment. Then using the m -convexity of S , clearly $S \sim \{p\}$ has at most $m - 1$ components S_1, \dots, S_k . Furthermore, it is easy to show that each set $\text{cl } S_i$ has at most $m_i - 1$ visually independent points and is m_i -convex, where $2 \leq m_i \leq m$ and where $\sum_{i=1}^k (m_i - 1) = m - 1$. Certainly $N \cap \text{cl } S_i \sim \{p\}$ is convex, so p cannot be an lnc point for any $\text{cl } S_i$. If we are able to show that each set $\text{cl } S_i$ is decomposable into $m_i - 1$ convex sets, then S will be

a union of $\sum_{i=1}^k (m_i - 1) = m - 1$ convex sets, finishing the argument. Hence it suffices to assume that $p \notin Q$, and the proof of Lemma 2 is complete.

Now repeat the construction used in the proof of Lemma 1 to define the collection \mathcal{R} of rays. Since Q is finite, we may order the rays in a clockwise direction, letting W_i denote the closed subset of S determined by consecutive rays R_i and R_{i+1} , $1 \leq i \leq n$, where $R_1 \cup R_{n+1} = H$. By previous remarks, we may assume that $W_i = \text{cl}(\text{int } W_i)$, $1 \leq i \leq n$. Thus to each i , $2 \leq i \leq n - 1$, there correspond two lnc points of S , denoted q_i, q'_i , where $q'_i = q_{i+1}$ for $1 \leq i \leq n - 1$. It is easy to show that each W_i set is convex, and we call W_i a *wedge* of S .

Again by the Lawrence, Hare, Kenelly Theorem, we may assume that bdy W_i is polygonal and select segments $[q_i, p_i]$, $[p'_i, q'_i]$ in bdy W_i , $1 \leq i \leq n$ (where $q_1 \in R_1, q'_n \in R_{n+1}$ are selected in the manner indicated previously).

We decompose S by defining $\mathcal{U}_1, \dots, \mathcal{U}_{m-1}$, each an appropriate collection of wedges of S . We assign wedges to the \mathcal{U}_i sets in the following manner: Let W_1 be in \mathcal{U}_1, W_2 in \mathcal{U}_2 , and let $P_1 = \{\mathcal{U}_1\}, P_2 = \{\mathcal{U}_1, \mathcal{U}_2\}$. Inductively, assume that each of the wedges W_1, \dots, W_j has been assigned to one of the sets $\mathcal{U}_1, \dots, \mathcal{U}_i$, and that $P_j = \{\mathcal{U}_1, \dots, \mathcal{U}_i\}$ partitions these j wedges so that $\text{conv}(\cup \mathcal{U}_i) \subseteq S, 1 \leq i \leq l$. Let V_i denote the last wedge assigned to \mathcal{U}_i (i.e., the wedge assigned to \mathcal{U}_i having largest subscript). If necessary, relabel the V_i and corresponding \mathcal{U}_i sets so that for $1 \leq i_1 < i_2 \leq l, V_{i_1}$ precedes V_{i_2} in our ordering. We assign W_{j+1} in the following manner: If $\text{conv}[W_{j+1} \cup (\cup \mathcal{U}_i)] \subseteq S$ for some i , choose i_0 to be the largest such subscript i , and assign W_{j+1} to \mathcal{U}_{i_0} . In this case, let $P_{j+1} = \{\mathcal{U}_1, \dots, \mathcal{U}_i\}$. If no such i exists, assign W_{j+1} to \mathcal{U}_{i+1} , and let $P_{j+1} = \{\mathcal{U}_1, \dots, \mathcal{U}_{i+1}\}$. In either case, P_{j+1} partitions the family $\{W_1, \dots, W_{j+1}\}$. Since there are finitely many wedges, the inductive procedure must end in a finite number of steps, and we may assume that the last partition $P_n = \{\mathcal{U}_1, \dots, \mathcal{U}_k\}$ partitions the family $\{W_1, \dots, W_n\}$ so that $\text{conv}(\cup \mathcal{U}_i) \subseteq S, 1 \leq i \leq k$. The integer k defined above will be called the *convex cover order* of S with respect to p . We will prove that $k = m - 1$.

For the remainder of the argument, we will let V_i denote the last wedge assigned to $\mathcal{U}_i, 1 \leq i \leq k - 1$, and let V_k denote the first wedge assigned to \mathcal{U}_k . Moreover, we relabel the \mathcal{U}_i and corresponding V_i sets so that for $1 \leq i_1 < i_2 \leq k, V_{i_1}$ precedes V_{i_2} in our clockwise ordering.

LEMMA 3. *In our assignment of wedges, for $1 \leq i < k, \text{conv}(V_i \cup V_k) \not\subseteq S$.*

PROOF OF LEMMA 3. Since $\text{conv}(\cup \mathcal{U}_i) \subseteq S$, it is easy to see that

$\text{conv}(V_k \cup (\cup \mathcal{U}_i))$ would lie in S , and in our partitioning procedure, V_k would have been assigned to \mathcal{U}_j for some $1 \leq j < k$, a contradiction.

LEMMA 4. *If S has convex cover order k with respect to p , then S contains k visually independent points x_1, \dots, x_k with $x_i \in (\text{int } U_i \sim \text{conv}(p, q_i, q'_i)) \cup (q_i, q'_i)$ for appropriate wedges U_1, \dots, U_k of S (ordered in a clockwise direction), $U_1 = V_1, U_k = V_k$. Moreover, x_1, x_k may be selected as close as we wish to q'_1, q'_k respectively.*

PROOF OF LEMMA 4. Applying induction, the assertion is trivial for $k = 1, 2$, so we assume it true for all positive integers less than k , to prove for k . Consider the collections of wedges $\mathcal{U}_1, \dots, \mathcal{U}_k$. By Lemma 3, $\text{conv}(V_1 \cup V_k) \not\subseteq S$, and there are two cases arising from the different ways $\text{conv}(V_1 \cup V_k)$ can contain points outside S : (1) For some wedge between V_1 and V_k , there is a corresponding member of Q which lies on the p side of $L(q'_1, q_k)$. (2) Case 1 does not occur and one of p'_1, p_k lies beyond $L(q'_1, q_k)$ from p .

Case 1. Suppose that for some wedge between V_1 and V_k , a corresponding member of Q lies on the p side of $L(q'_1, q_k) = L$. Among the q and q' points having this property, examine those whose distance to L is maximal, and from these select the one having largest subscript j . Assume it is $q'_{j-1} = q_j$ and corresponds to W_j .

Let M be the line through q_j parallel to L . Now the set S' defined as the union of V_1, W_{j-1} , and the wedges of S following V_1 and preceding W_{j-1} is clearly of the type considered in the hypothesis and has convex cover order $t < k$ according to the procedure involved in our assignment of wedges. (Clearly t is the largest integer l for which members of \mathcal{U}_l are contained in S' .) Note that if $V_t \subseteq S'$, then V_t must be the last wedge in S' , since every wedge following V_t has been assigned to \mathcal{U}_j for some $j > t$ (by our labeling). By the induction hypothesis, S' contains points x_1, \dots, x_t visually independent via S' (and hence via S), and satisfying our specified requirements. Similarly let S'' be the union of V_k and the wedges of S following V_t and preceding V_k . Then S'' satisfies our hypothesis. Also, each of V_{t+1}, \dots, V_k lies in S'' , so S'' has convex cover order at least $k - t$. Hence there are points y_{t+1}, \dots, y_k in S'' which are visually independent via S'' (and via S) and which satisfy our requirements. Certainly $\text{int}(S' \cap S'') = \emptyset$, each of x_1, \dots, x_t precedes the ray $R(p, q_j)$, and each of y_{t+1}, \dots, y_k follows $R(p, q_j)$.

We assert that $x_1, \dots, x_t, y_{t+1}, \dots, y_k$ are visually independent via S : Clearly by our selection of q_j , for $1 < r \leq t$, each x_r is in $S' \sim V_1$ and must lie on M or

beyond M from p . For $t + 1 \leq s < k$, each y_s is in $S'' \sim V_k$ and must lie beyond M from p . Also, by our induction hypothesis, we may assume that x_t and y_k are beyond M . Therefore, if $[x_r, y_s] \subseteq S$ for some $1 \leq r \leq t$ and some $t + 1 \leq s \leq k$, then $q_i \in \text{int conv}\{p, x_r, y_s\} \subseteq \text{int } S$, and q_i could not be an lnc point for S , clearly impossible. Hence the points are indeed visually independent. The remaining part of the inductive hypothesis is easy to verify, and the proof of Case 1 is complete.

Case 2. Suppose that Case 1 does not occur and that p_k lies beyond $L(q'_1, q_k) = L$ from p . (The proof for p'_i beyond L is similar and will be omitted.) Let S' denote the union of V_1 and the wedges of S which follow V_1 and precede V_k . Then S' has convex cover order $k - 1$ and we may apply our induction hypothesis to obtain visually independent points x_1, \dots, x_{k-1} . Since every lnc point between q'_1 and q_k is on or beyond L , our inductive hypothesis assures us that x_2, \dots, x_{k-1} see no points on $[p_k, q_k)$ via S . Again by our hypothesis we may assume that x_1 sees no point on $[p_k, q_k)$ either. Then for an appropriate choice of x_k in V_k , x_1, \dots, x_k are visually independent via S . The rest of the argument is easy, finishing Case 2 and the proof of Lemma 4.

Now since S is m -convex, S has no more than m visually independent points. Thus by Lemma 4, S must have convex cover order $m - 1$, and S is the union of the $m - 1$ convex sets $S_i = \text{conv}(\cup \mathcal{U}_i)$, $1 \leq i \leq m - 1$, completing the proof of the theorem.

COROLLARY 1. *Let S be a closed set in the plane. Let $p \in \ker S \neq \emptyset$, and for H some line containing p , assume $S \subseteq \text{cl } H_1$. Then for $m \geq 2$, S is m -convex if and only if S is the union of $m - 1$ convex sets.*

COROLLARY 2. *Let S be a closed set in the plane, Q the set of lnc points of S , with $p \in [(\text{bdry } S) \cap (\ker S)] \sim Q \neq \emptyset$. Then for $m \geq 2$, S is m -convex if and only if S is the union of $m - 1$ convex sets.*

PROOF. Since $p \notin Q$, we may select a neighborhood N of p such that $N \cap S$ is convex. Then a hyperplane H supporting $N \cap S$ at p will have the required property.

COROLLARY 3. *Let S be a closed m -convex set in the plane with $p \in \ker S \neq \emptyset$. Then S is a union of $2(m - 1) = 2m - 2$ or fewer convex sets.*

PROOF. Let H be any line through p and apply Theorem 1 to $\text{cl } H_1 \cap S$ and $\text{cl } H_2 \cap S$.

3. Applications of Theorem 1

In this section we present three remarkably different kinds of decomposition theorems which may be proved from Theorem 1.

THEOREM 2. *Let S be a closed m -convex set in R^2 with $p \in \ker S \cap \text{bdry } S \neq \emptyset$. Then S is the union of $2m - 3$ or fewer convex sets.*

PROOF. The set S may obviously be expressed as the union of an m -convex set and an $(m - 1)$ -convex set, each satisfying the hypothesis of Theorem 1.

THEOREM 3. *If S is any closed 4-convex subset of the plane, then S is the union of 6 or fewer convex sets.*

PROOF. If S is not simply connected, then S is the union of 5 or fewer convex sets by a theorem of Guay [3]. Hence assume S is simply connected. Also, we may assume that S is connected, for otherwise the bound may be lowered to 4. In case $Q = \emptyset$, S is convex [7], and the result is trivial. For $Q \neq \emptyset$, select q in Q and define sets $S_q = \{x : [x, q] \subseteq S\}$ and $S^q = \{x : [x, q] \not\subseteq S\}$ (called the star and anti-star of q in S , respectively). The set S_q is closed and since $q \in \text{bdry } S$, we have $q \in \ker S_q \cap \text{bdry } S_q$. Using the simple connectedness of S , it is easy to show that S_q is 4-convex, and by Theorem 2, S_q is the union of $2 \cdot 4 - 3 = 5$ convex sets C_1, \dots, C_5 . Using the fact that $q \in Q$, it is easy to show that for x, y in S^q , $[x, y] \subseteq S$, and again by the simple connectedness of S , $C_6 = \text{conv } S^q \subseteq S$. Hence $S = \cup_{i=1}^6 C_i$, the desired result.

REMARK. By Example 3 in [5], the bound for a closed planar 4-convex set is no lower than 5. Hence the best bound is either 5 or 6.

The final two theorems of this section deal with the case in which $\text{conv } Q \subseteq S$.

THEOREM 4. *Let S be a closed m -convex set in the plane, $m \geq 2$, with Q the set of lnc points of S . If $\text{conv } Q \subseteq S$ and $\text{int conv } Q = \emptyset$, then S is expressible as a union of $m - 1$ convex sets. The bound is best possible for every m .*

PROOF. By comments in [1], we may assume that no component of $S \sim Q$ is a segment. Also, we may assume that Q is not a singleton set, for then the proof is easy. If S is 2-convex the result is trivial, and we assume the result is true when S is j -convex, $2 \leq j < m$, to prove for m .

In case $S \sim Q$ is not connected, then $S \sim Q$ has at most $m - 1$ components S_1, \dots, S_k , $k \leq m - 1$. It is not hard to show that each set $\text{cl } S_i$ has at most $m_i - 1$ visually independent points and is m_i -convex, where $2 \leq m_i < m$ and where

$\Sigma_{i=1}^k(m_i - 1) = m - 1$. Then by our induction hypothesis each set $\text{cl } S_i$ is a union of $m_i - 1$ convex sets, and S is a union of $\Sigma_{i=1}^k(m_i - 1) = m - 1$ convex sets, finishing the argument.

In case $S \sim Q$ is connected, then it is easy to show that S has $m - 2$ lnc points, and S is a union of $m - 1$ convex sets by a theorem of Guay and Kay [4, Theorem 1]. Clearly the bound of $m - 1$ is best possible, and the proof is complete.

THEOREM 5. *Let S be a closed m -convex set in the plane, $m \geq 2$, with Q the set of lnc points of S . If $\text{conv } Q \subseteq S$, then S is a union of $3m - 2$ or fewer convex sets.*

PROOF. The proof of this result is lengthy. First we shall show we may assume that Q is finite (Lemma 5), next that each component of $S \sim \text{conv } Q$ is convex (Lemma 6), and that $S \sim Q$ is connected (Lemma 7). Then we define a subset T of S satisfying Theorem 1, Corollary 2, and show that the remaining points of S lie either in $\text{conv } Q$ or in one of at most $2(m - 1)$ components of $S \sim \text{conv } Q$. For the sake of brevity, some of the easy details of the argument are omitted.

Without loss of generality, assume S is connected, for every component of S not containing Q is necessarily convex. Also, it is clear that S must be simply connected. By earlier remarks, we may assume that S is bounded and that $S = \text{cl}(\text{int } S)$. And by Theorem 4, we may restrict our attention to the case in which $\text{int } \text{conv } Q \neq \emptyset$.

LEMMA 5. *To any finite subset $F = \{x_i : 1 \leq i \leq k\}$ of S there corresponds an m -convex set T having finitely many lnc. points, with $F \subseteq T \subseteq S$. Hence we may assume that S has finitely many lnc points.*

PROOF OF LEMMA 5. Let $F = \{x_i : 1 \leq i \leq k\}$ be any finite subset of S , and without loss of generality, assume that the points of F are indexed so that $x_i \in S \sim \text{conv } Q$ for $1 \leq i \leq n$ and $x_i \in \text{conv } Q$ for $n \leq i \leq k$. By a lemma of Valentine [10, Lemma 1], each point in F sees some point of Q (and hence some point of $\text{conv } Q$) via S . Moreover, if $x \notin \text{conv } Q$, then x necessarily sees some point y in $\text{bdry } \text{conv } Q$ such that $[x, y) \cap \text{conv } Q = \emptyset$. Therefore, for each i , $1 \leq i \leq n$, we may define $A_{x_i} = A_i = \{y : y \text{ in } \text{bdry } \text{conv } Q, [x_i, y) \subseteq S \text{ and } [x_i, y) \cap \text{conv } Q = \emptyset\}$. Also, since S is simply connected and $\text{conv } Q \subseteq S$, A_i is necessarily connected, so A_i is an arc in $\text{bdry } \text{conv } Q$.

We assert that the endpoints v_i, w_i (not necessarily distinct) of the arc A_i lie in Q : Let $a_i, b_i \in A_i$ and assume for the moment that a_i, b_i may be selected so that $a_i \neq b_i$. Then no point on A_i between a_i and b_i may lie in Q . In case $a_i \in Q$, then no point in $\text{bdry } \text{conv } Q$ beyond $L(x_i, a_i)$ from b_i may lie in A_i , and we may select

$v_i = a_i$. If $a_i \notin Q$, then it is not hard to show that there are points of Q beyond $L(x_i, a_i)$ from b_i , and we may select such a point v so that the arc $\widehat{va_i}$ in $\text{bdry conv } Q$ has minimal length. Then $[v, a_i] \cup [a_i, x_i] \subseteq S$, no point of Q is in $\text{conv } \{v, a_i, x_i\} \sim [v, x_i]$, so by a result of Valentine [8, Corollary 2], $\text{conv } \{v, a_i, x_i\} \subseteq S$, and x_i sees v via S . Then using the fact that $\widehat{va_i}$ has minimal length, it is easy to show that $v \in A_i$. Since $v \in Q$, x_i can see no point of $\text{bdry conv } Q$ beyond $L(x_i, v)$ from a_i , and $v = v_i$ is the required point. A similar argument holds for b_i to produce w_i , finishing the argument. In case $\{a_i\} = \{b_i\} = A_i$, the previous argument may be adapted to show that $a_i \in Q$, and the assertion is proved.

For each $i, 1 \leq i \leq n$, let W_i denote the component of $S \sim \text{conv } Q$ containing x_i , and let B_i denote the subset of $\text{conv } Q$ corresponding to W_i — i.e., $B_i = \text{cl } W_i \cap \text{bdry conv } Q$. By earlier remarks, it is clear that for y in $W_i, A_y \subseteq B_i$. Moreover, since S is locally starshaped [5, Lemma 2], for s in B_i , there is some y in W_i such that $s \in A_y$, and $B_i = \cup \{A_y : y \text{ in } W_i\}$. Now for s, t in B_i , we may select y, z in W_i such that $[y, s] \cup [z, t] \subseteq S$. Since W_i is locally convex and connected, it is polygonally connected, and there is a path λ in $W_i \subseteq S \sim \text{conv } Q$ from s to t . Then λ cannot intersect $[s, t] \subseteq \text{conv } Q$, and it is easy to show that there is an arc in B_i from s to t . Clearly B_i is closed, so B_i is an arc $\widehat{q_i q'_i}$. Also, by earlier remarks, $q_i, q'_i \in Q$.

Define $T = (\cup_{i=1}^n W_i) \cup (\text{conv } Q)$. It is easy to show that T is an m -convex subset of S . Moreover, since S is locally starshaped and W_i is polygonally connected, an earlier argument may be adapted to show that the set of lnc points of T lies in $\{q_i, q'_i : \widehat{q_i q'_i} = \text{cl } W_i \cap \text{bdry conv } Q, 1 \leq i \leq n\}$. Therefore, by the Lawrence, Hare, Kenelly Theorem, it suffices to assume that Q is finite, finishing the proof of Lemma 5.

LEMMA 6. *Without loss of generality, we may assume that each component W_i of $S \sim \text{conv } Q$ is convex.*

PROOF OF LEMMA 6. Assume that some component W_i of $S \sim \text{conv } Q$ is not convex, and let $Q_i \neq \emptyset$ denote the set of lnc points of $\text{cl } W_i$. By the proof of Lemma 5, $Q_i \subseteq \{q_i, q'_i\}$, where $\widehat{q_i q'_i} = \text{cl } W_i \cap \text{bdry conv } Q$. Let L be a line which contains q_i, q'_i and which supports $\text{conv } Q$, and let L_1, L_2 denote the corresponding open halfspaces, with $\text{conv } Q \subseteq \text{cl } L_2$. Then using an argument employed in [4, lemma 6], $L_1 \cap W_i$ is convex and each of the two (or fewer) components of $L_2 \cap W_i$ is convex. If $L_2 \cap W_i$ has two components, then S may be written as the union of two convex sets and an $(m - 2)$ -convex set T . In case $L_2 \cap W_i$ is connected, then S is the union of a convex set and an $(m - 1)$ -convex set T . In

either case, $T \cap W_i$ is convex, and without loss of generality we may assume that W_i is convex.

LEMMA 7. *Without loss of generality we may assume that $S \sim Q$ is connected.*

PROOF OF LEMMA 7. If $S \sim Q$ is not connected, then it is easy to show that S is expressible as the union of a convex set $\text{cl } W_i$ and an $(m - 1)$ -convex set $\text{cl}(S \sim W_i)$, for W_i a component of $S \sim Q$.

Returning to the proof of the theorem, order the lnc points of S and the corresponding components of $S \sim \text{conv } Q$ in a clockwise direction along $\text{bdry conv } Q$. By Lemma 7, to each component W_i of $S \sim \text{conv } Q$ there corresponds a pair q_i, q'_i of lnc points of S (where q'_i follows q_i in our ordering). By the Lawrence, Hare, Kenelly Theorem, we may assume that $\text{bdry } W_i$ is polygonal, and hence we may select segments $[q_i, p_i], [p'_i, q'_i]$ in $\text{bdry } W_i$. Let $L_i = L(q_i, p_i)$, $L'_i = L(q'_i, p'_i)$. We will say that a point x is *beneath* L_i if x is in the open halfspace L_{i1} determined by L_i and containing W_i . Similarly, x is *beyond* L_i if x is in the open halfspace L_{i2} .

Let W be any fixed component in $S \sim \text{conv } Q$, and for convenience, assume $W = W_1$. We assert that $[q_1, q'_1]$ fails to be in $\text{cl}(L_{i1})$ for at most $m - 1$ of the lines L_i , $1 < i \leq n$: Assume that $[q_1, q'_1] \not\subseteq \text{cl}(L_{i1}) \cup \text{cl}(L_{j1})$, where $1 < i < j \leq n$. Then p_j sees no point of S beyond L_j . But q'_1 is necessarily beyond L_j , and hence q_1 is, too (since if $q_1 \neq q'_1$, then q_1 follows q'_1 and precedes q_j in our ordering). Thus for c_i selected appropriately in (p_i, q_i) , $[c_i, q_i]$ lies beyond L_j and no point of $[c_i, q_i]$ sees any point of $[p_j, q_j)$ via S . Then clearly for every collection of halfspaces $\text{cl}(L_{i1})$ which fail to contain $[q_1, q'_1]$, there is a corresponding collection of visually independent points of S , so at most $m - 1$ halfspaces have this property. Let \mathcal{A} denote the associated collection of components of $S \sim \text{conv } Q$.

Similarly, letting \mathcal{B} denote the collection of components W_i of $S \sim \text{conv } Q$ for which $[q_1, q'_1] \not\subseteq \text{cl}(L'_{i1})$, then \mathcal{B} has at most $m - 1$ members. Define $T = \text{cl}(S \sim \cup \{W : W = W_1, W \in \mathcal{A} \text{ or } W \in \mathcal{B}\})$. For $W \notin \mathcal{A} \cup \mathcal{B}$, every point of W sees $[q_1, q'_1]$ via T , and for t in (q_1, q'_1) , t is $\text{bdry } T \cap \text{ker } T$. Since T is a closed m -convex set, we may use Theorem 1, Corollary 2, to conclude that T is a union of $m - 1$ convex sets. Therefore, S is a union of $3(m - 1) + 1 = 3m - 2$ or fewer convex sets, finishing the proof of Theorem 5.

4. The general case

A general decomposition theorem will require several preliminary lemmas.

LEMMA 8. *Let S be a closed m -convex set in the plane. If B is the closure of a*

bounded component A of $R^2 \sim S$, then $\text{conv } B$ is a polygon having at most $m - 1$ sides.

PROOF OF LEMMA 8. Certainly $\text{conv } B$ is the convex hull of its extreme points. To see that $\text{conv } B$ is a polygon, we show that it has at most $2m - 1$ (and hence finitely many) extreme points: If $\text{conv } B$ had $2m$ extreme points, they could be ordered in a clockwise direction along $\text{bdry conv } B$. Letting x_1, x_2, \dots, x_{2m} denote these points, clearly the set $\{x_{2k} : 1 \leq k \leq m\}$ would be a set of m visually independent points of S , for otherwise A could not be connected. However, this would contradict the m -convexity of S . Thus $\text{conv } B$ may have at most $2m - 1$ extreme points.

It remains to show that the polygon $\text{conv } B$ has at most $m - 1$ sides. Let x_1, x_2, \dots, x_k denote the vertices of $\text{conv } B$, $k \geq 3$, where the points are again ordered in a clockwise direction along $\text{bdry conv } B$. Then $A \subseteq B \subseteq \text{conv } B$. We will select k visually independent points y_1, \dots, y_k of S . (For convenience of notation, let $x_{k+1} = x_1$.) If (x_i, x_{i+1}) contains a point in S , let $y_i \in (x_i, x_{i+1}) \cap S$. Otherwise, (x_i, x_{i+1}) lies in a (possibly unbounded) component of $R^2 \sim S$, and this component is distinct from A since $A \subseteq \text{conv } B$. Hence it is not hard to see that there is some component S_i of $S \sim \{x_i, x_{i+1}\}$ which lies in $\text{conv } B$ such that $x_i, x_{i+1} \in \text{cl } S_i$. In this case, select $y_i \in S_i$. Then y_1, \dots, y_k is a visually independent subset of S , for otherwise A could not be connected. Therefore $k \leq m - 1$ and the proof of Lemma 8 is complete.

LEMMA 9. *Let $S = \text{clint } S$ be a set in the plane. If $R^2 \sim S$ has at least $r = (n + 2)2^{n-1}$ bounded components having closures B_1, \dots, B_n , $n \geq 0$, and for each i , $\text{conv } B_i$ is a convex polygon, then S has at least $n + 3$ visually independent points on $\bigcup_{i=1}^n \text{bdry } B_i$.*

PROOF OF LEMMA 9. The proof is by induction. If $n = 0$, then $r = 1$ and clearly S has at least three visually independent points. Assume the result is true for integers less than n , $n \geq 1$, to prove for n .

Consider the polygon $P \equiv \text{conv}(\bigcup_{i=1}^n B_i)$ and let p be any extreme point of P . Then p is an extreme point of some $\text{conv } B_i$, say of $\text{conv } B_r$. Choose a line H supporting P such that $H \cap P = \{p\}$. Choose a line L through p which intersects $\text{int } B_r$. Now if $n + 2$ of the sets B_1, \dots, B_r share a segment with L , then we can find $n + 3$ visually independent points in $\bigcup_{i=1}^n \text{bdry } B_i$, finishing the argument. For since $S = \text{clint } S$, no two B sets share a segment. To each B_i sharing a segment with L , we may associate points p_i, p'_i on L with $B_i \cap L$ containing $[p_i, p'_i]$ and p_i, p'_i in $\text{bdry } B_i$. Also, we may relabel the B sets and corresponding p

points so that $p_1 < p'_1 \leq p_2 < \dots \leq p_{n+2} < p'_{n+2}$. Clearly by selecting points x_i in $\text{bdry } B_i \sim L$, x_i sufficiently close to p_i , and y_{n+2} in $\text{bdry } B_{n+2} \sim L$, y_{n+2} sufficiently close to p'_{n+2} , we have $x_1, \dots, x_{n+2}, y_{n+2}$ a set of $n + 3$ visually independent points of S .

Hence we may assume that L meets at most $n + 1$ of the sets B_1, \dots, B_r in a segment. Then from these r sets there are at least

$$r - (n + 1) = (n + 2)2^{n-1} - n - 1$$

not sharing a segment with L , and each of these sets must lie entirely in one of the closed halfspaces determined by L . Hence one of these halfspaces, say $\text{cl}(L_1)$, must contain at least $(r - n - 1)/2$ of the B_i sets, $1 \leq i \leq r - n - 1$.

Consider the set

$$S' = S \cup (\cup \{B_i : 1 \leq i \leq r, B_i \not\subseteq \text{cl}(L_1)\}),$$

Then S' is a closed set having at least $r' \geq (r - n - 1)/2$ bounded components which satisfy the hypothesis of the theorem. Moreover

$$\begin{aligned} \frac{r - n - 1}{2} &= \frac{(n + 2)2^{n-1} - n - 1}{2} \\ &= (n + 1)2^{n-2} + 2^{n-2} - \frac{n}{2} - \frac{1}{2}. \end{aligned}$$

If $n \geq 3$, then $2^{n-2} \geq n/2 + 1/2$, and thus $r' \geq (r - n - 1)/2 \geq (n + 1)2^{n-2}$. In case $n = 1$, then $(r - n - 1)/2 = (3 - 2)/2 = 1/2$, and since r' is an integer, $r' \geq 1 = (n + 1)2^{n-2}$. Similarly, if $n = 2$ then $(r - n - 1)/2 = (4 \cdot 2 - 3)/2 = 5/2$, and $r' \geq 3 = (n + 1)2^{n-2}$. We conclude that for $n \geq 1$, S' has at least $(n + 1)2^{n-2}$ bounded components in its complement. Therefore, by the induction hypothesis, S' has at least $n + 2$ visually independent points x_1, \dots, x_{n+2} on $\cup \{\text{bdry } B_i : 1 \leq i \leq r\} \cap \text{cl}(L_1)$, and these points are also visually independent via S .

We assert that we may choose a point x_{n+3} on $\text{bdry } B_r$ which sees none of the points x_1, \dots, x_{n+2} via S : Recall that the line L contains points interior to B_r . Let q denote the vertex of the polygon $\text{conv } B_r$ which lies in the open halfspace L_2 and for which $[p, q] \subseteq \text{bdry conv } B_r$. In case $(p, q) \cap \text{bdry } B_r \neq \emptyset$, then let x_{n+3} be any member of this set. Otherwise, fix $y \in (p, q)$ and consider the set $\{x : x \in \text{bdry } B_r \text{ and } (y, x) \cap B_r \neq \emptyset\}$. Then let x_{n+3} be any member of this set in L_2 and distinct from q . Clearly x_{n+3} sees no point of S in $\text{cl}(L_1)$, and x_1, \dots, x_{n+3} is a set of visually independent points. This completes the induction and finishes the proof of Lemma 9.

COROLLARY 1. *If $S = \text{clint } S$ is an m -convex set in the plane, then $R^2 \sim S$ has no more than $(m - 1)2^{m-4} - 1$ bounded components.*

PROOF. By Lemma 8, if B is the closure of a bounded component of $R^2 \sim S$, then $\text{conv } B$ is a polygon (having at most $m - 1 \geq 3$ edges). Then by Lemma 9, if $R^2 \sim S$ has r bounded components, assuming $(n + 2)2^{n-1} \leq r < (n + 3)2^n$, $n \geq 0$, then S has at least $n + 3$ visually independent points. Since $n + 3 \leq m - 1$, we have $r < (m - 1)2^{m-4}$, the desired result.

LEMMA 10. *Let $S = \text{clint } S$ be an m -convex set in the plane, $m \geq 3$, with $x \in S$. Then the set $S_x = \{y : [x, y] \subseteq S\}$ is k -convex, where $2 \leq k \leq (m - 1)2^{m-4} + 1$.*

PROOF OF LEMMA 10. Suppose on the contrary that S_x contains at least $k = (m - 1)2^{m-4} + 1 \geq m$ visually independent points x_1, \dots, x_k . Let A_1, \dots, A_r denote the bounded components of $R^2 \sim S$, where $0 \leq r \leq (m - 1)2^{m-4} - 1$ by the corollary to Lemma 9. For each nonempty set A_i , select a point b_i in A_i and examine the rays $R(x, b_i)$. Order the rays in a clockwise direction and relabel the A_i sets so that $R(x, b_{i+1})$ follows $R(x, b_i)$ in our ordering.

The rays define closed subset of the plane: If $r \geq 2$, let T_i be the closed subset determined by $R(x, b_i)$ and $R(x, b_{i+1})$ relative to our clockwise orientation, $1 \leq i \leq r$ (where $r + 1 \equiv 1$). At most one T_i set is not convex, and if this occurs, bisect the corresponding angle to yield two convex sets, T_{i1} and T_{i2} . Hence we obtain either r or $r + 1$ closed convex T sets. In case $r = 1$, let T_1, T_2 be the closed halfspaces determined by the line $L(x, b_1)$, and if $r = 0$, let T_1 be the plane. Clearly in all cases $T_i \cap S$ is simply connected for each i . (Otherwise some ray $R(x, b_i)$ would lie between $R(x, b_i)$ and $R(x, b_{i+1})$ in our ordering, impossible.)

We assert that at least m of the points x_1, \dots, x_k lie in one of the convex T regions: If fewer than m of the points x_1, \dots, x_k belonged to T_i for each i , then there would be at most $(m - 1)(r + 1) \leq (m - 1)2^{m-4} < k$ points in all, a contradiction.

Hence one of the regions, say T_1 , contains m of the x points. For convenience, say $x_1, \dots, x_m \in T_1$. By m -convexity of S , at least one corresponding segment, say $[x_1, x_2]$, is in S . Hence $[x, x_1] \cup [x, x_2] \cup [x_1, x_2] \subseteq S$ with $[x_1, x_2] \not\subseteq S_x$, denying simple-connectedness of $T_1 \cap S$. Thus S_x is indeed k -convex, $2 \leq k \leq (m - 1)2^{m-4} + 1$, finishing the proof of Lemma 10.

THEOREM 6. *If S is any closed m -convex set in the plane, $m \geq 3$, then S is the union of $(m - 1)2^{m-3}$ or fewer convex sets.*

PROOF. Without loss of generality, we may assume that $S = \text{clint } S$, for otherwise S is the union of k segments and an $(m - k)$ -convex set for some $1 \leq k \leq m - 2$.

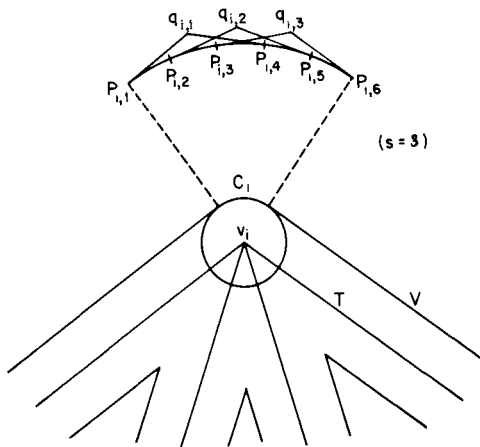
By [5, Theorem 2], there exist $m - 1$ or fewer points x_1, \dots, x_{m-1} in S such that $S = \bigcup_{i=1}^{m-1} S_i$, where $S_i = \{y : [x_i, y] \subseteq S\}$. By Lemma 10, each S_i is at most $[(m - 1)^2 2^{m-4} + 1]$ -convex. Since $x_i \in \ker S_i \neq \emptyset$, by Corollary 3 to Theorem 1, each S_i is a union of $2[(m - 1)^2 2^{m-4}]$ or fewer convex sets. Thus S is the union of $(m - 1)[2(m - 1)^2 2^{m-4}] = (m - 1)^3 2^{m-3}$ or fewer convex sets.

COROLLARY 1. A closed set S in the plane is m -convex for some $m \geq 2$ if and only if S is the union of finitely many closed convex sets.

5. An example

M. A. Perles has communicated the following example of a class of compact m -convex subsets of E^2 , mentioned in an earlier paper of one of the authors [5].

Example. For each integer $r \geq 2$ and $s \geq 1$ a set $S_{r,s}$ will be defined by first taking the vertices v_1, \dots, v_r of a regular polygon in E^2 inscribed in a unit circle, and setting $T = \cup \{[v_i, v_j] \mid 1 \leq i < j \leq r\}$. With $0 < \delta < \pi/10r$ put $V = T + \delta B$, where B is the unit disk. Hence, V consists of $\binom{r}{2}$ parallel strips of width 2δ joining one another at the points v_i , with the outer corners being rounded off by disks of radius δ centered at the v_i , $i = 1, \dots, r$. If K is the boundary of the set $\text{conv } V$ then K consists of segments parallel to $[v_i, v_{i+1}]$ and circular arcs C_i of radius δ , $i = 1, \dots, r$, where each C_i is less than a semicircle. Now divide each C_i into $2s - 1$ equal sub arcs with consecutive points of division labeled $p_{i,1}, \dots, p_{i,2s}$ (see figure). For each pair $(p_{i,j}, p_{i,s+j})$ let the tangents to C_i at $p_{i,j}$ and $p_{i,s+j}$ meet at



$q_{i,j}$, $j = 1, \dots, s$, and put $\Delta_{i,j} = \text{conv}\{p_{i,j}, p_{i,s+j}, q_{i,j}\}$. Finally, define $S_{r,s} = V \cup (\cup \{\Delta_{i,j} \mid 1 \leq i \leq r, 1 \leq j \leq s\})$.

It can be easily verified that $S_{r,s}$ has the following properties, for each $r \geq 2$, $s \geq 1$:

- 1) The points $q_{i,1}, \dots, q_{i,s}$ are visually independent via $S_{r,s}$.
- 2) One may associate a point with each of the $\binom{r}{2}$ parallel strips so that the resulting $\binom{r}{2}$ points are visually independent via $S_{r,s}$.
- 3) Starting with the s points $q_{i,1}, \dots, q_{i,s}$ for any i , points may be associated with each of the remaining $\binom{r-1}{2}$ parallel strips not passing through v_i yielding $s + \binom{r-1}{2}$ visually independent points.

4) $S_{r,s}$ is m -convex, where $m = 1 + \max\left\{\binom{r}{2}, s + \binom{r-1}{2}\right\}$.

5) Each of the points $q_{i,k}$ can see each of $q_{j,l}$ via $S_{r,s}$ for each $i \neq j$, $1 \leq k \leq s$, $1 \leq l \leq s$; consequently, $\text{conv}(\Delta_{i,k} \cup \Delta_{j,l}) \subset S_{r,s}$ for $1 \leq i \leq r$, $1 \leq j \leq r$, $1 \leq k \leq s$, $1 \leq l \leq s$, and $k \neq l$.

6) If r is even and $s \geq r$, in order to cover $S_{r,s}$ with the least number of convex subsets, choose the $\binom{r}{2}$ parallel strips together with one $\Delta_{i,j}$ at each end per strip, leaving $s - r + 1$ sets $\Delta_{i,j}$ not accounted for at each v_i . Opposite pairs of these remaining $\Delta_{i,j}$ can be taken into convex subsets inside $r/2$ parallel strips yielding

$$\binom{r}{2} + (s - r + 1)\frac{r}{2} = \frac{rs}{2} = \left\lceil \frac{rs + 1}{2} \right\rceil$$

convex sets. A similar analysis yields the same number when r is odd.

7) If $s < r$ then all the sets $\Delta_{i,j}$ can be included with the $\binom{r}{2}$ parallel strips.

8) Since $\left\lceil \frac{rs + 1}{2} \right\rceil \leq \binom{r}{2}$ if and only if $s \leq r - 1$, $S_{r,s}$ is the union of n closed, convex sets and is not the union of fewer than n , where

$$n = \max\left\{\binom{r}{2}, \left\lceil \frac{rs + 1}{2} \right\rceil\right\}.$$

Note that if $s < r$ then $S_{r,s}$ is an example of an m -convex set which is the union of $m - 1$ but no fewer convex sets, since in this case $m - 1 = \binom{r}{2} = n$. But

consider the set $S = S_{r,s}$, where $s = r^2 \geq r$ ($r \geq 2$). S is then a compact, planar m -convex set with

$$m = 1 + s + \binom{r-1}{2} = \frac{3r^2 - 3r + 4}{2}, \quad n = \left\lceil \frac{rs + 1}{2} \right\rceil \cong \frac{r^3}{2}.$$

It then follows that

$$\sqrt{\frac{2}{27}}m^{3/2} < n < \sqrt{\frac{2}{9}}m^{3/2},$$

so S is the union of less than $m^{3/2}$ convex sets but is not the union of $(1/4)m^{3/2}$ convex sets, and m can assume the values of the infinite sequence 5, 11, 20, 32, \dots .

Hence, the best possible bound, while possibly not as large as $(m-1)^{3m-3}$, cannot be linear in m , in general. In higher dimensions, the situation is infinitely worse since Perles has also constructed an example of a compact 3-convex subset of E^4 which is not the union of a finite number of convex sets.

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